#### NATIONAL UNIVERSITY OF SINGAPORE

### CS4226 - INTERNET ARCHITECTURE

(Semester 1: AY2016/17)

Time Allowed: 2 Hours

#### **INSTRUCTIONS TO STUDENTS**

- 1. Please write your Student Number only. Do not write your name.
- 2. This assessment paper contains SIX (6) questions and comprises TWELVE (12) printed pages.
- 3. Students are required to answer ALL questions.
- 4. All questions must be answered in the space provided in the answer sheet; no extra sheets will be accepted as answers.
- 5. This is a CLOSED BOOK assessment.
- 6. You are allowed to bring one A4 help sheet. No book is allowed.
- 7. Electronic calculators are not allowed.

STUDENT	NO:			

This portion is for examiner's use only

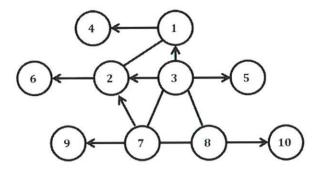
Question	Marks	Remarks
Q1		
Q2		
Q3	N 200-200	
Q4		
Q5		
Q6		
Total		

# Question 1: Exponential Distribution and M/M/1 System [15 marks]

tem at a random to t an event could b	time, what is the expose either an arrival of	pected waiting t r a departure. E	ime until you see xplain your deriv	the next evation.	vent? Notice [10 marks]

## **Question 2: ISP Peering Relationship [20 marks]**

Two common ISP peering relationships are (1) provider-customer and (2) peer-peer relationships. In each case, the bilateral business agreements will disallow certain paths. In the following figure, each node represents an ISP, each directed edge represents a provider-customer relationship (arrow side is the customer, for example, ISP 4 is a customer of ISP 1), and each undirected edge represents a peer-peer relationship.



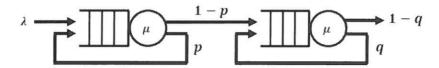
**A.** For each of the following paths, indicate whether it is valid based on the ISPs' business relationships. [12 marks]

path	is it valid?	path	is it valid?
$4 \rightarrow 1 \rightarrow 2 \rightarrow 7 \rightarrow 9$		$10 \rightarrow 8 \rightarrow 7 \rightarrow 2 \rightarrow 1 \rightarrow 4$	
$4 \rightarrow 1 \rightarrow 2 \rightarrow 6$		$1 \rightarrow 3 \rightarrow 7 \rightarrow 8 \rightarrow 10$	
$9 \rightarrow 7 \rightarrow 3 \rightarrow 8 \rightarrow 10$		$5 \rightarrow 3 \rightarrow 2 \rightarrow 7 \rightarrow 9$	

**B.** BGP policy routing can be used to control traffic. Give examples of how an ISP could use BGP attributes to control its incoming and outgoing traffic. Please state what you are trying to achieve, and how that is achieved by using the specific BGP attribute. [8 marks]

Description of the Control of the Co		

## **Question 3: Jackson Network [15 marks]**

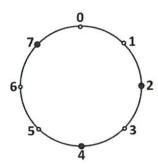


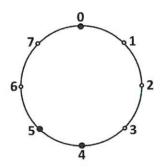
In the above Jackson network, we assume  $0 and the arrival rate <math>\lambda$  is chosen so as to not overload the network.

Determine the average sojourn time of jobs $\mathbf{E}[W]$ as a function of $\lambda, \mu, p$ and $q$ . [10 marks]

<b>D.</b> If we interchange the order of the queues (that is $p$ and $q$ are flipped), does in	$\mathbb{E}[W]$ increase,
decrease, or stay the same? Please justify your answers.	[5 marks]

# Question 4: Peer-to-Peer Networks [18 marks]





**A.** Consider a Chord network with namespace [0,7]. Suppose three nodes 2, 4 and 7 are active as shown in the above left figure, please construct the finger tables for the three active nodes. [9 marks]

start	interval	successor

 ger table o interval	successor

start	interval	successor

**B.** Suppose after nodes 0 and 5 join and nodes 2 and 7 leave the network, three nodes 0, 4 and 5 are active as shown in the above right figure. Please construct the finger tables for the three active nodes. [9 marks]

fin	ger table o	of node 0
start	interval	successor
-11 - 11		

start	interval	successor	

start	interval	successor	
-			

# **Question 5: Fair Resource Allocation [16 marks]**

Consider a network with four links 1,2,3 and 4 that have capacities $C_1 = 1$ , $C_2 = C_4 = 2$ (Mbps), respectively. There are four traffic flows: flow 1 traverses links 1 traverses links 2 and 3; flow 3 traverses links 3 and 4; and flow 4 traverses links 4 at the demand of the four flows are $d_1 = 2$ , $d_2 = 3$ , $d_3 = 4$ and $d_4 = 1$ (Mbps), respectively. Calculate the max-min fair allocation $\mathbf{x} = (x_1, x_2, x_3, x_4)$ to the four flows.	and 2; flow 2 nd 1. Suppose
<b>B.</b> Identify the bottleneck link(s) for each traffic flow under the max-min allocation	on. [4 marks]

Calculate the weighted max-min fair allocation $\mathbf{x} = (x_1, x_2, x_3, x_4)$ to weights are $\phi = (\phi_1, \phi_2, \phi_3, \phi_4) = (1, 2, 3, 4)$ .	the four flows when [4 marks]
Identify the bottleneck link(s) for each traffic flow under the weighter Part C.	ed max-min allocation [4 marks

## **Question 6: Scheduling [16 marks]**

Consider a single router with two packet flows A and B. The router has a processing capacity of 50 bytes/second. Before (clock) time t = 0, the router is empty. The first two packets from flow A arrive at time  $t_a^1 = 1$  (seconds) and  $t_a^2 = 2$  (seconds) with length  $t_a^1 = t_a^2 = 35$  (bytes). The first two packets from flow B arrive at time  $t_b^1 = 0$  (seconds) and  $t_b^2 = 3$  (seconds) with length  $t_b^1 = t_b^2 = 90$  (bytes). No other packets arrive afterwards.

 $\mathbf{A}$ . If the two flows have the weights  $4\phi_A = \phi_B = 4$ , calculate the real (or wall clock) finishing time  $f_a^1, f_a^2, f_b^1$  and  $f_b^2$  for each packet under GPS. [5 marks]

р г		e the virtual finishing ti	[6 marks

The two flows have the weights $4\phi_A = \phi_B$ me $\hat{F}_a^1$ , $\hat{F}_a^2$ , $\hat{F}_b^1$ and $\hat{F}_b^2$ for each packet under WF me, we use FIFO to break the tie.	FQ. If two packets have t	he same virtual finishing [5 marks]

Scratch Paper

— END OF PAPER —