

**NATIONAL UNIVERSITY OF SINGAPORE****CS4226 – INTERNET ARCHITECTURE**

(Semester 1: AY2018/19)

Time Allowed : 2 Hours

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**INSTRUCTIONS TO STUDENTS**

1. Please write your Student Number only. Do not write your name.
2. This assessment paper contains **SIX (6)** questions and comprises **TEN (10)** printed pages.
3. Students are required to answer **ALL** questions.
4. All questions must be answered in the space provided in the answer sheet; no extra sheets will be accepted as answers.
5. This is a CLOSED BOOK assessment.
6. You are allowed to bring one A4 help sheet. No book is allowed.

STUDENT NO: \_\_\_\_\_

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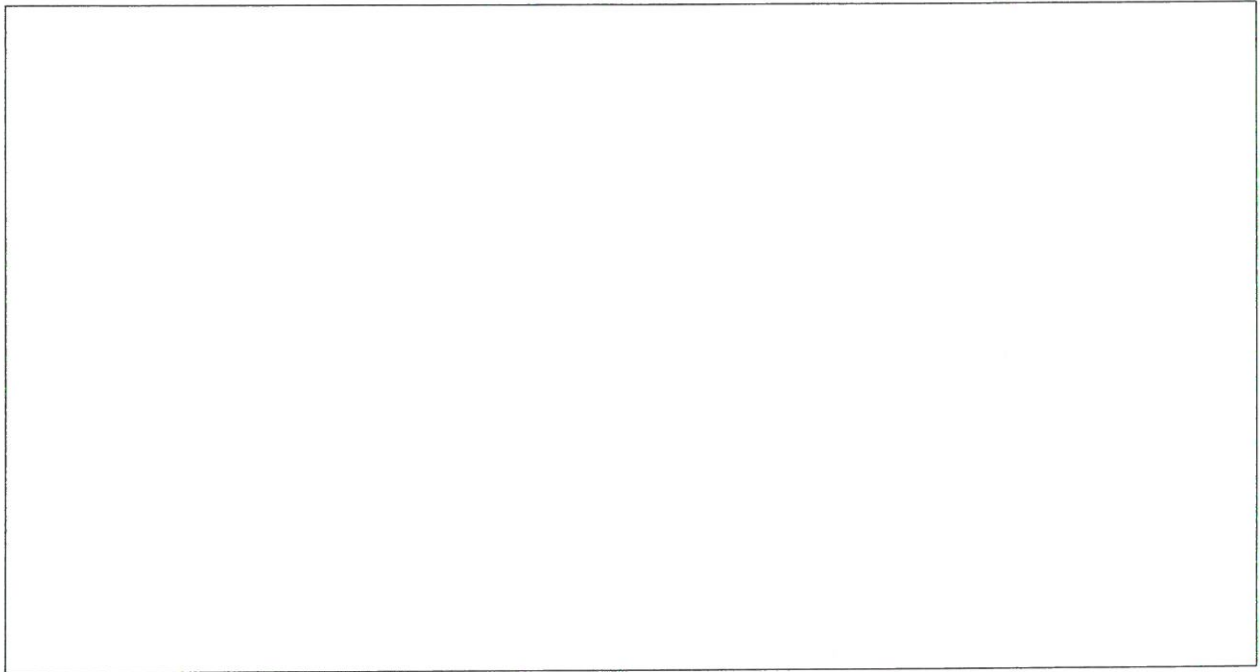
This portion is for examiner's use only

Question	Marks	Remarks
Q1		
Q2		
Q3		
Q4		
Q5		
Q6		
Total		

**Question 1: M/M/1 Link Model [10 marks]**

Consider a link modeled as an  $M/M/1$  system. Suppose we know that the mean sojourn time of the packets is  $E[W] = 3$  seconds and 80% of time that the server is busy.

**A.** What is the packet flow's arrival rate to the link? [5 marks]

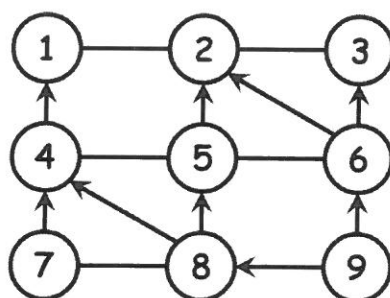


**B.** What is the minimal arrival rate of the packet flow that is going to make the system unstable? [5 marks]



### Question 2: BGP and Inter-Domain Routing [20 marks]

Two common ISP peering relationships are (1) provider-customer and (2) peer-peer relationships. In each case, the bilateral business agreements will disallow certain paths. In the following figure, each node represents an ISP, each directed edge represents a provider-customer relationship (arrow side is the customer, for example, ISP 4 is a customer of ISP 1), and each undirected edge represents a peer-peer relationship.

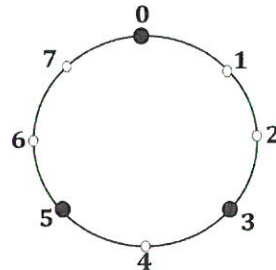


- A. For each of the following paths, indicate whether it is valid based on the ISPs' business relationships. [12 marks]

path	is it valid?	path	is it valid?
4 → 1 → 2 → 3 → 6		9 → 8 → 5 → 2 → 6	
4 → 1 → 2 → 5 → 8		1 → 4 → 7 → 8	
7 → 4 → 5 → 8		3 → 2 → 5 → 8 → 9	

- B. Describe the differences between iBGP and eBGP in terms of their functionalities. [8 marks]

**Question 3: Peer-to-Peer Networks [20 marks]**



Consider a Chord network with namespace  $[0, 7]$  and three nodes 0, 3 and 5 being active.

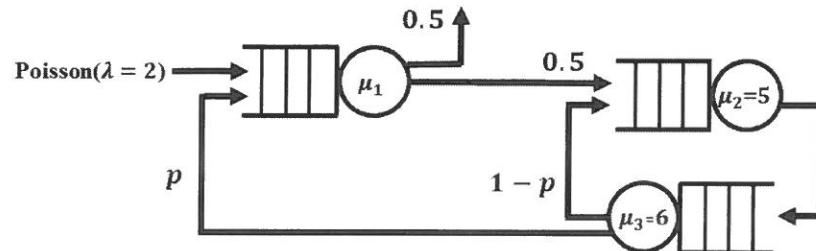
**A.** Please construct the finger tables for node 0 and node 3. [10 marks]

finger table of node 0		
start	interval	successor

finger table of node 3		
start	interval	successor

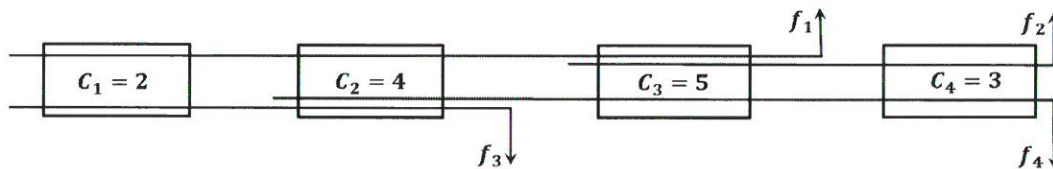
**B.** Suppose some content with a hash key value of 0 is queried at node 3. Show the sequence of nodes visited until the content is retrieved. [5 marks]

**C.** Suppose some content with a hash key value of 2 is queried at node 5. Show the sequence of nodes visited until the content is retrieved. [5 marks]

**Question 4: Jackson Network [10 marks]**

In the above Jackson network, the service rates of the right two servers are  $\mu_2 = 5$  (packets/second) and  $\mu_3 = 6$  (packets/second). The arrival rate is  $\lambda = 2$  (packets/second). To guarantee the stability of the system, what are the conditions we need for the routing probability  $p$  and the server capacity  $\mu_1$ ? [10 marks]

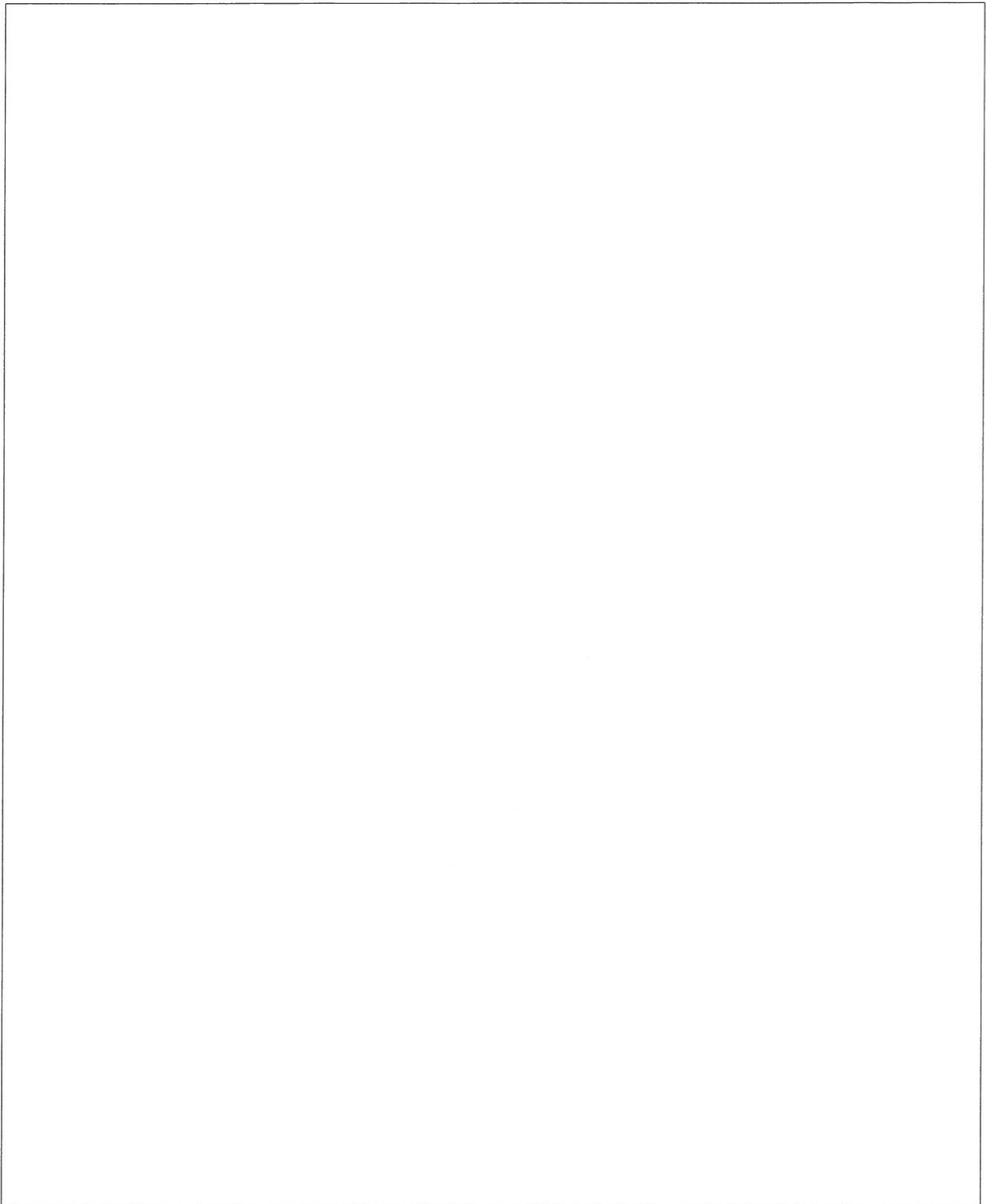
### Question 5: Fair Resource Allocation and Scheduling [20 marks]



Consider a network path with four links 1, 2, 3 and 4 that have capacities  $C_1 = 2$ ,  $C_2 = 4$ ,  $C_3 = 5$  and  $C_4 = 3$  (Mbps), respectively. There are four traffic flows: flow  $f_1$  traverses links 1, 2 and 3; flow  $f_2$  traverses links 3 and 4; flow  $f_3$  traverses links 1 and 2; flow  $f_4$  traverses links 2, 3 and 4. Suppose the demand of the four flows are  $d_1 = 2$ ,  $d_2 = 4$ ,  $d_3 = 1$  and  $d_4 = 3$  (Mbps), respectively.

**A.** Calculate the weighted max-min fair allocation  $\mathbf{x} = (x_1, x_2, x_3, x_4)$  to the three flows when the weights are  $\phi = (\phi_1, \phi_2, \phi_3, \phi_4) = (1, 2, 3, 4)$ . [10 marks]

**B.** Identify the bottleneck link(s) for each traffic flow under the above weighted max-min allocation. [10 marks]



**Question 6: Weighted Fair Queueing (WFQ) [20 marks]**

Consider a single switch serving two packet flows  $a$  and  $b$ . The switch has a processing capacity of 80 bytes/second. Before (clock) time  $t = 0$ , the switch is empty. The first two packets from flow  $a$  arrive at time  $t_a^1 = 0$  (seconds),  $t_a^2 = 3$  (seconds) with length  $l_a^1 = l_a^2 = 120$  (bytes). The first two packets from flow  $b$  arrive at time  $t_b^1 = 1$  (seconds) and  $t_b^2 = 2$  (seconds) with length  $l_b^1 = 60$  (bytes) and  $l_b^2 = 180$  (bytes). No other packets arrive afterwards.

**A.** If the two flows have the weights  $3\phi_a = \phi_b = 3$ , calculate the real (or wall clock) finishing time  $f_a^1, f_a^2, f_b^1$  and  $f_b^2$  for each packet under GPS. [5 marks]



**B.** If the two flows have the weights  $3\phi_a = \phi_b = 3$ , calculate the virtual finishing time  $F_a^1, F_a^2, F_b^1$  and  $F_b^2$  for each packet under WFQ. [10 marks]

**C.** If the two flows have the weights  $3\phi_a = \phi_b = 3$ , calculate the real (or wall clock) finishing time  $\hat{F}_a^1, \hat{F}_a^2, \hat{F}_b^1$  and  $\hat{F}_b^2$  for each packet under WFQ. If two packets have the same virtual finishing time, we use FIFO to break the tie. [5 marks]

Scratch Paper

— END OF PAPER —